Cuttlefish: A Fair, Predictable Execution Environment for Cloud-hosted Financial Exchange

Liangcheng (LC) Yu, Prateesh Goyal, Ilias Marinos, and Vincent Liu

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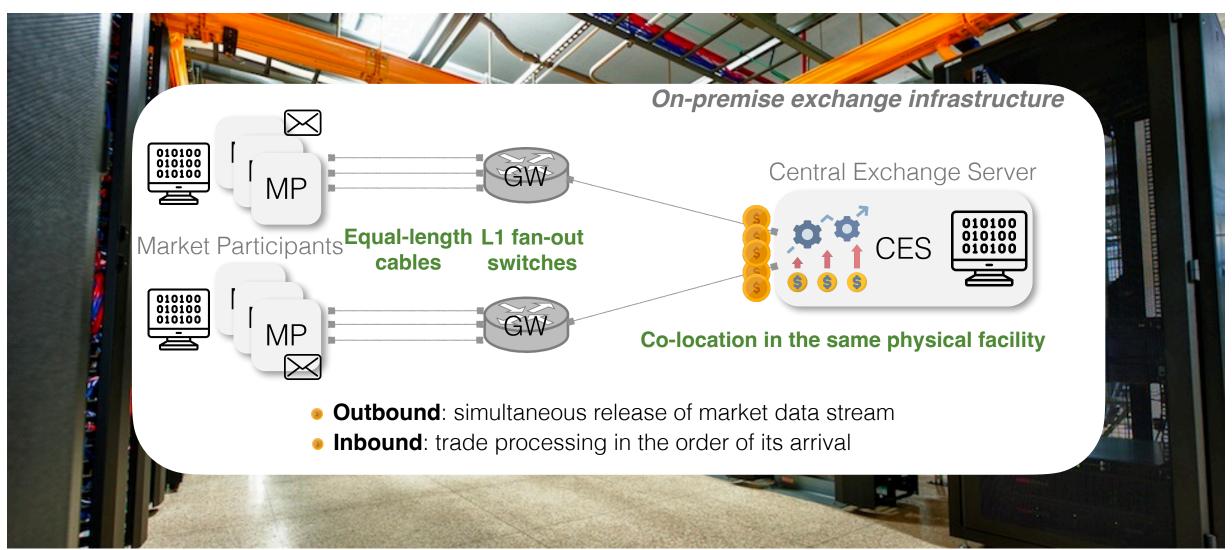








On-premise exchange infrastructure



Interest in cloud-hosted exchange services



- System scalability and resource elasticity
- Rise of remote work
- Cost reduction and ease of management

. . .

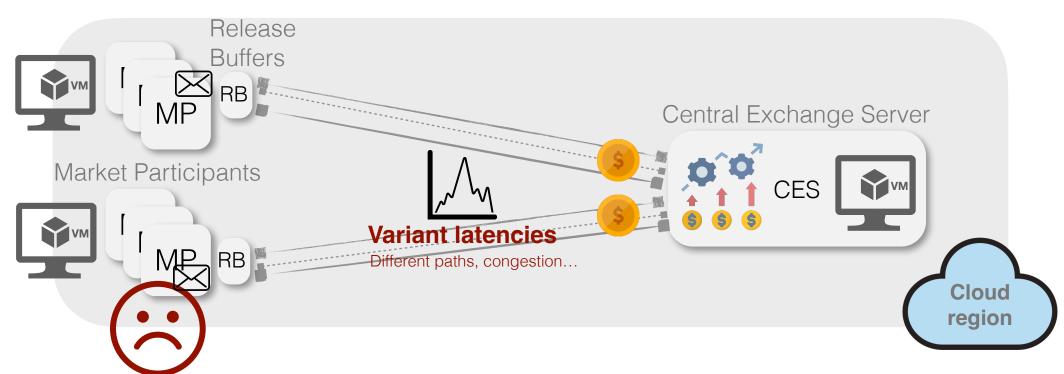
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Variances in network latencies

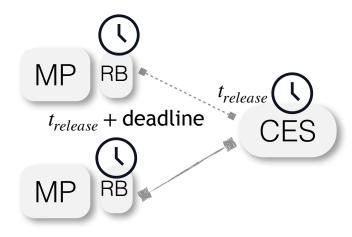


- Outbound: simultaneous release of market data stream
- Inbound: trade processing in the order of its arrival

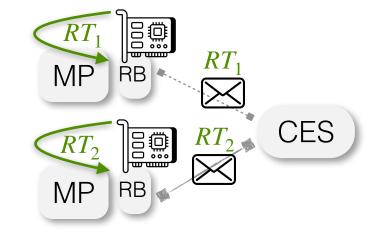
Unfairness!

Efforts toward communication fairness

Clock synchronization (CloudEx, HotOS '21)



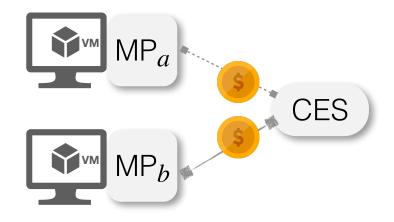
Logical clock based on response time (RT) (DBO, SIGCOMM '23)



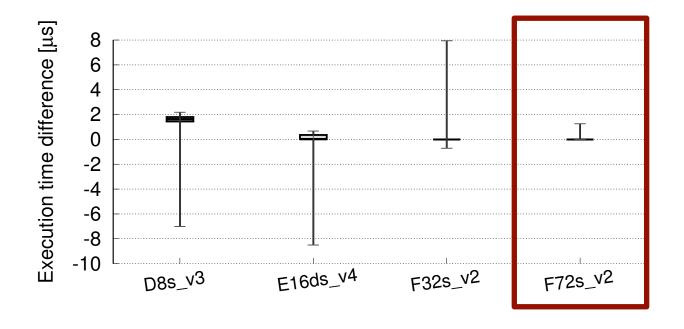
- Perfect clock synchronization is **hard**
- Hard to pre-determine the deadline
- **Example 2** Limited to trigger-point based trades
- **Doesn't handle MP-RB latency variances**

...cloud execution can also incur unfairness!

Other sources of unfairness: noisy neighbors, thermal conditions of the processors...

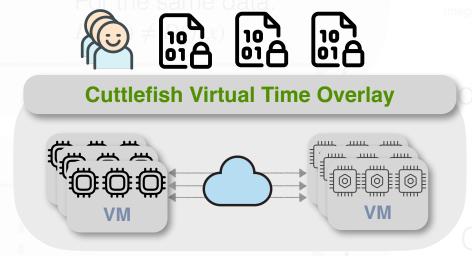


Identical programs running on same types of VMs



Can we **eliminate variations** that come from the cloud infrastructure?







Equal cloud networks (thermal condition, resource

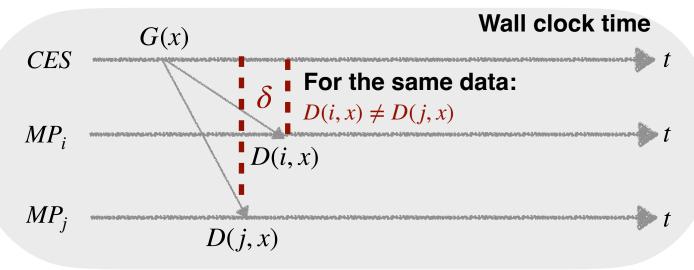
Equal execution hardware

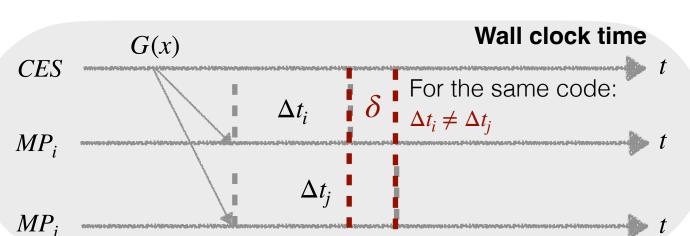
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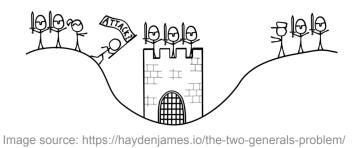
Outline

- Conceptual foundation
- Implementing virtual time overlay
- Evaluation

Let's reflect on underlying model today...



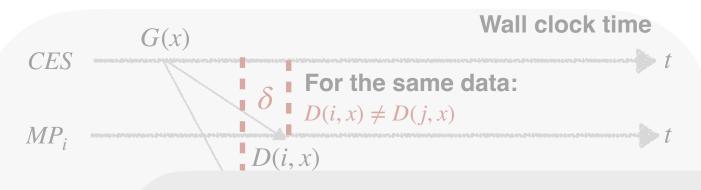


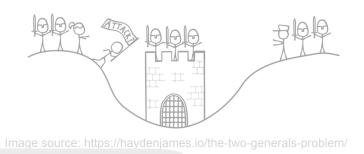


Simultaneous delivery is *infeasible*!

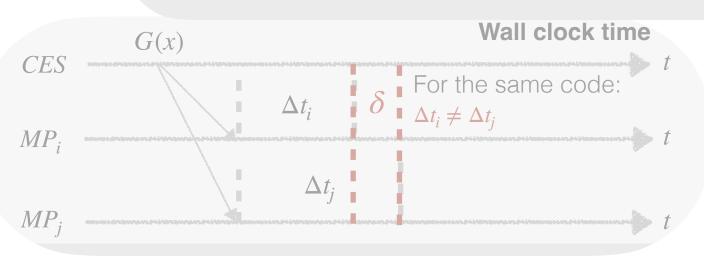
Execution time can be **non-deterministic** at $O(\mu s)$ (thermal condition, resource utilization...)

Let's reflect on underlying model today...





Communication and computation synchrony are *challenging* in real-time



Execution time can be **non-deterministic** at $O(\mu s)$ (thermal condition, resource utilization...)

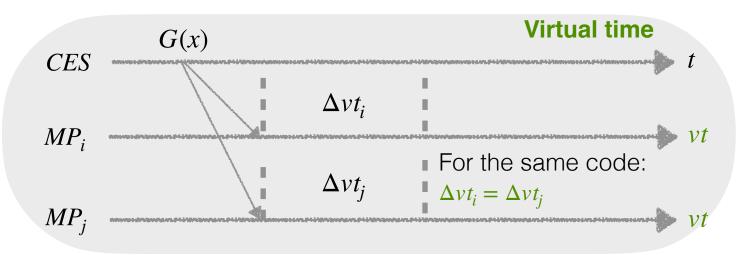
delivery

ible!



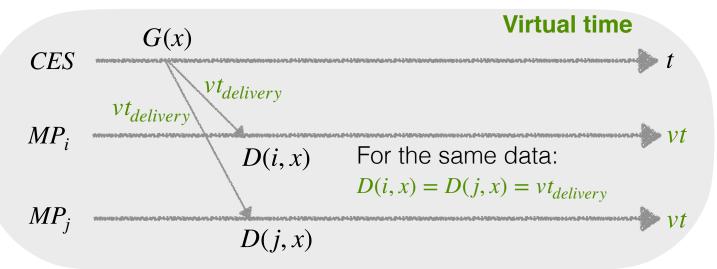
Let's try virtual time domain ...

Virtual time unit ≡ some equal amount of work



Execution synchrony:

Advancing virtual time per 'actual amount of work'



Communication synchrony:

Releasing data to MPs at the same virtual delivery time

Outline

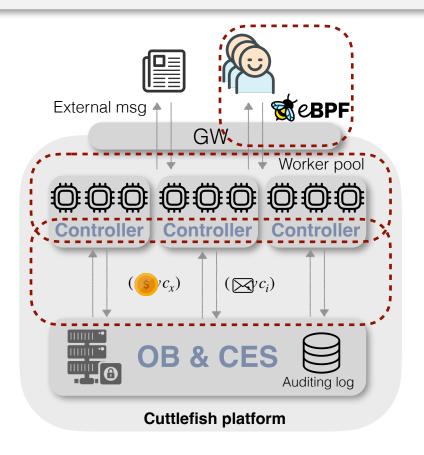
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Implementing virtual time abstraction



Instantiate vt as virtual cycles of a platform-agnostic IR/VM

Account and control the advancement of virtual cycles

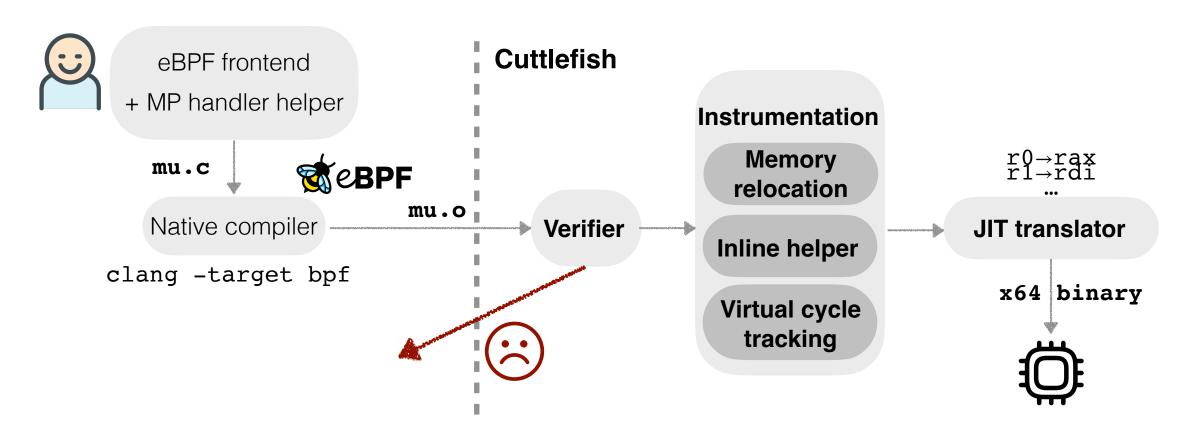


- Programming interface
- Runtime execution
- Virtual cycle tracking

User programming abstraction

```
White-list set of
#include <cuttlefish_user.h>
                                        extensible service APIs
int mp handler(subscribed context t* data) {
    if ((*data) > 100) {
        // Sell
        trade t trade = 1;
         submit trade(&trade);
    } else if ((*data) < 10) {</pre>
        // Buy
                                           Just-in-time trade
        trade t trade = 2;
                                              submission
        submit trade(&trade);
    map update(0, &trade);
                                   Narrow KV store API (e.g., lookup,
    return 0;
                                     update) for stateful invocations
```

MP code lifetime

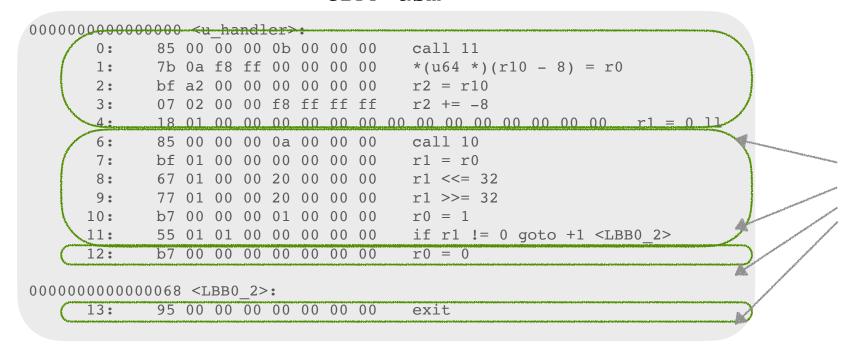


2-tier compilation with the platform agnostic IR:

Track virtual cycle (fairly) in eBPF, but *execute* (efficiently) on native HW target

How to track and advance virtual time cycles?

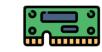
eBPF asm



Native HW asm

; movabs r11, <vc address>
49 BB F0 DE BC 9A 78 56 34 12
; add qword ptr [r11], 2
49 81 03 02 00 00 00
x64

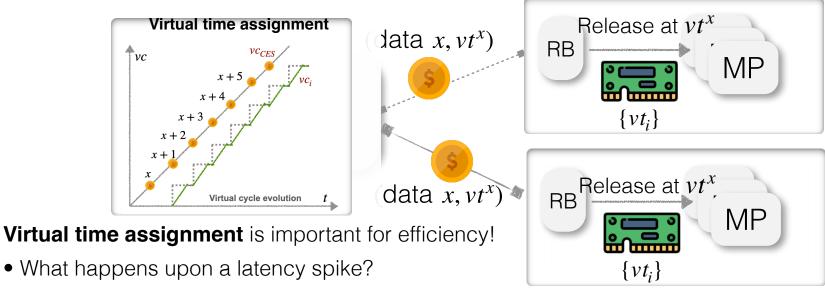
 $vt_i + = \Delta vt$



 $\{vt_i\}$ maintained by execution runtime

- ullet Break into basic blocks for batch updates of vt_i
 - JMP source, JMP destination, trade submission call
- Emit native machine code (2 x64 instr.) at the epilogue during JIT translation
 - Dummy trade/heartbeat for large blocks
 - Update the offsets for the (direct) JMP instructions

Simultaneous data delivery in virtual time

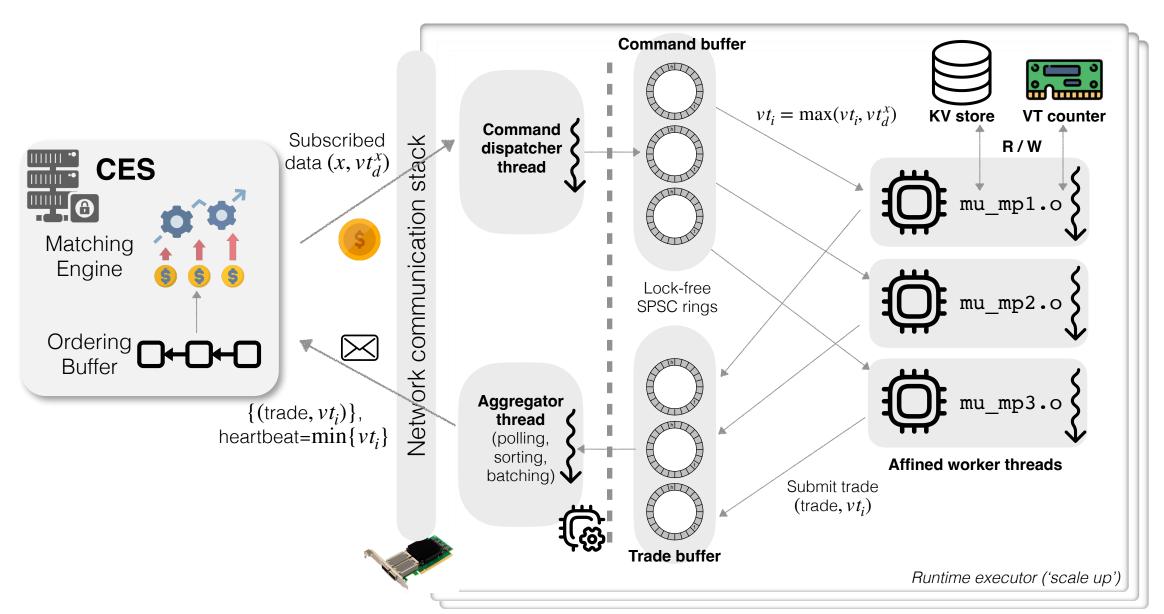


What if some processor executions get slower?

More details:

- Virtual time assignment algorithm
- Fault tolerance
- Handling external messages
- Security & trust discussions

Runtime execution workflow



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Comparison with existing ordering schemes

Ordering mechanisms

- Response Time (RT) based ordering
- FIFO ordering

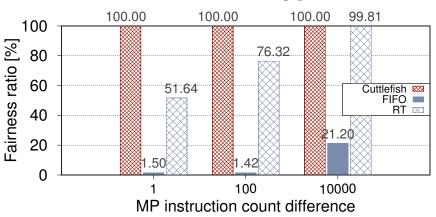
Set up

- Two MPs on two VMs
- ullet MP $_a$ executes N additional primitive IR instructions than MP $_b$
- Market data rate: every $\approx 100 \mu s$

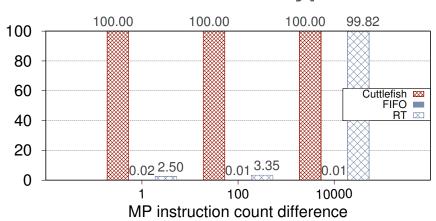
Metric

Fairness ratio

Same VM type



Different VM types



Performance cost for fairness

Set up

- 100 MPs on 10 VMs
- Market data rate: every $\approx 100 \mu s$
- CX-4 NIC and Intel Xeon Platinum 8272CL CPU @ 2.60GHz

	Latency (µs)				
	avg.	$\mathbf{p50}$	$\mathbf{p}90$	p 99	$\mathbf{p}99.9$
MaxRTT	52.04	47.74	49.95	55.85	144.2
Cuttlefish	54.19	50.82	53.49	68.46	166.3
	+2.15	+3.08	+3.54	+12.61	+22.1

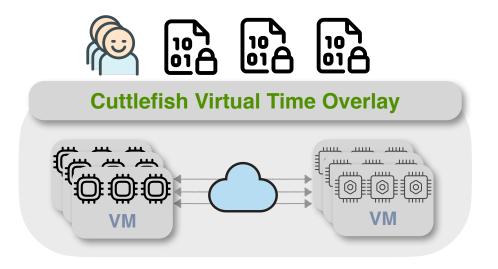
More details:

- Execution throughput and latencies under processor disparities
- Virtual time instrumentation overhead
- Recovery under failures

Summary

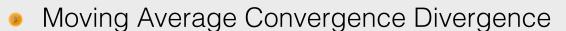
Cuttlefish: a fair, predictable cloud-hosted exchange platform

- Abstracting out variances in cloud communication and execution hardware
- An efficient implementation runnable on commercial cloud



The interface is expressive enough

- Fibonacci, Bubble Sort...
- SMA Mean Reversion
- EMA Mean Reversion
- Relative Strength Index
- Moving Average Crossover Strategy
- Bollinger Bands Strategy



- Multiple Moving Average Crossover Strategy
- Parabolic SAR
- On Balance Volume (OBV) + EMA
- Stochastic Oscillator
- Basic Market Making
- (§)



